

IOMax: Maximizing Out-of-Core I/O Analysis Performance on HPC Systems

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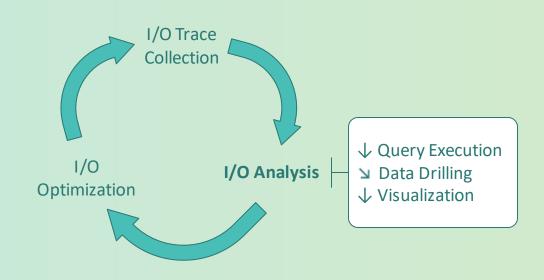
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This work was performed under the auspices of the U.S. Department of Energy by Lawrence Livermore National Laboratory under Contract DE-AC52-07NA27344. This material is based upon work supported by the U.S. Department of Energy, Office of Science, Office of Advanced Scientific Computing Research under the DOE Early Career Research Program (LLNL-PRES-857387). Also, the material is based upon work supported by the National Science Foundation under Grant no. NSF OAC-2104013, OCI-1835764, and CSR-1814872.

I/O Analysis Cycle

- Optimizing I/O efficiency has become essential for maximizing productivity as scientific workloads on HPC systems become increasingly data-intensive
- The go-to solution: I/O Analysis
 - Involves gathering I/O traces and examining patterns to detect anomalies
 - Existing I/O analysis tools utilize data drilling to identify I/O bottlenecks within trace data that can fit in memory
- Recently, traces from scientific workloads have reached terabytes in size, necessitating out-of-core analysis



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"data drilling": Iterative exploration of data for deeper insights



"out-of-core analysis":
Analysis of data too large to
fit in memory by means of
distributed computing

State-of-the-art

- Currently, I/O analysis is conducted using various tools, including
 - Profiling/Tracing: Darshan, Recorder, DLIO Profiler
 - Analysis: PyDarshan, Drishti, DXT Explorer, Recorder-viz
- Studies mostly rely on these tools to detect I/O problems, for instance
 - UMAMI [1], TOKIO [2], IOMiner [3]: System-wide, Darshan, Pandas, In-memory
 - Extracting and characterizing I/O behavior of HPC workloads [4]: Workflow-level, Recorder, Parquet, Out-of-core

Challenges

• I/O analysis on large-scale I/O traces, involving data drilling, is a complex task that faces three major challenges:



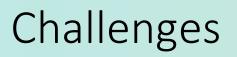
Challenges

24/2024

- As the volume of traces generated by scientific workloads has grown to terabytes in size, there is an increasing need for out-of-core I/O analysis
- For instance, traces from DLIO Benchmark [5] for a TensorFlow workload reaches 5TB in size
 - Accesses 5376 files of 132MB in size, 4 read threads, 1000 epoch
- Calculating the unoverlapped I/O requires efficient slicing and grouping in async I/O scenarios

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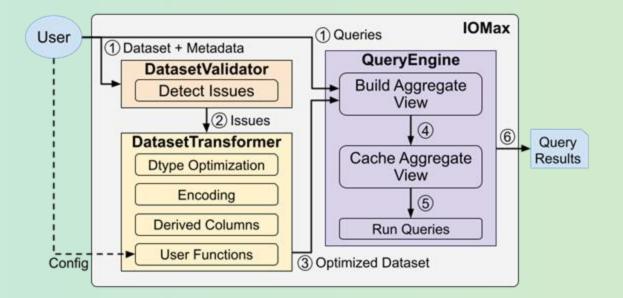
• Typical queries to analyze this trace would look like the following in SQL

<pre># Total I/O size by each thread within a specific time SELECT SUM(size) FROM traces WHERE time > '01:50' and time < '01:53' GROUP BY thread_id</pre>	
<pre># Aggregated I/O BW per process within a specific time SELECT SUM(size)/SUM(duration) FROM traces WHERE time > '01:50' and time < '01:53' GROUP BY process_id</pre>	

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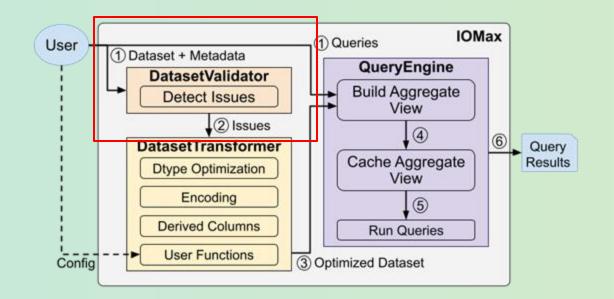
Solution: IOMax

- The IOMax tool is aimed at enhancing the efficiency of data drilling analysis on large-scale I/O traces
- It has three main components
 - DatasetValidator
 - DatasetTransformer
 - QueryEngine



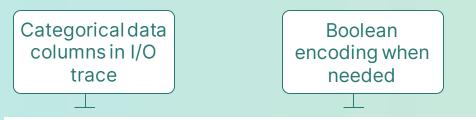
$IOMax \rightarrow DatasetValidator$

- Detects issues within the I/O traces
 - Misinferred datatypes
 - Inefficient datatypes
 - Encoding issues (binary/string)
- For instance
 - I/O Category and Access Pattern columns are among commonly misinferred datatypes
 - String columns are highly inefficient for analysis, both when storing and querying them

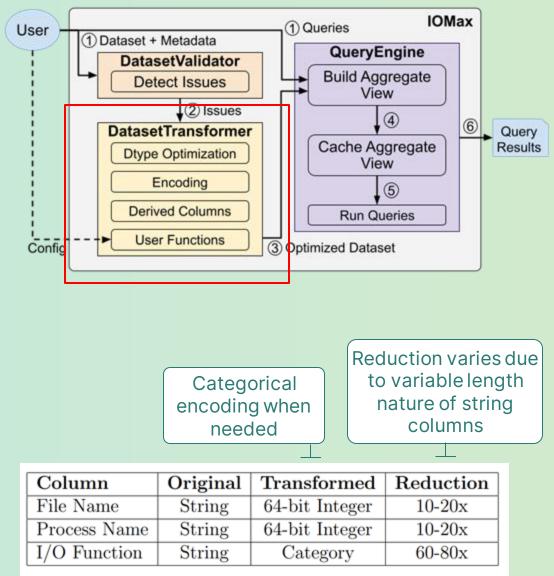


IOMax → DatasetTransformer

- Corrects misinferred datatypes by converting them into the expected datatypes
- For instance
 - I/O Category and Access Pattern columns usually have limited set of values, hence can be transformed into more efficient datatypes
 - Strings columns such as File and Process Names can be encoded as integers values through hashing

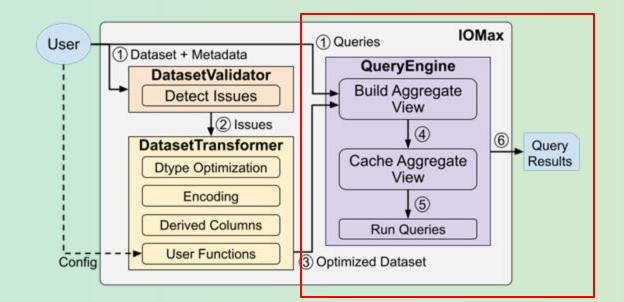


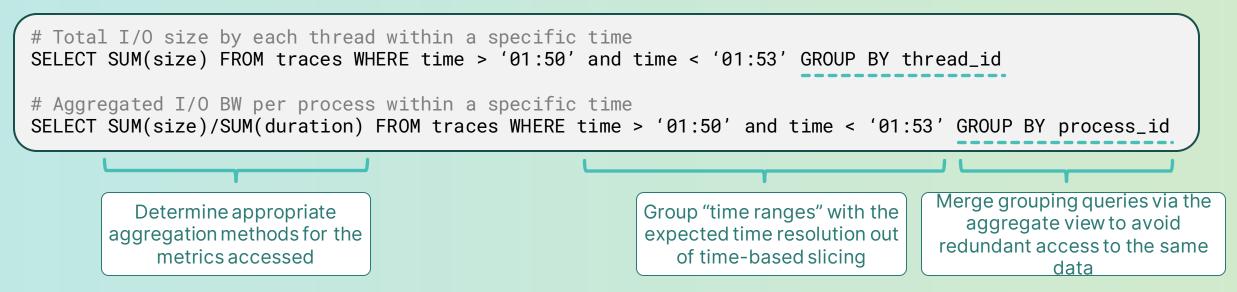
Column	Misinferred	Transformed	Reduction
I/O Category	32-bit Integer	8-bit Integer	$\sim 4x$
Access Pattern	32-bit Integer	Boolean	$\sim 4x$



$IOMax \rightarrow QueryEngine$

 Constructs an execution plan that leverages query reduction and in-memory caching techniques to minimize redundant I/O costs





Evaluations

- Workloads
 - Microbenchmarks utilizing varied I/O trace records
 - 5 million, 25 million, 125 million
 - Derived from real-world I/O traces
 - 4 scientific HPC workflows
 - 1000 Genomes, Montage (in this presentation)
 - HACC, CM1 (in the paper)

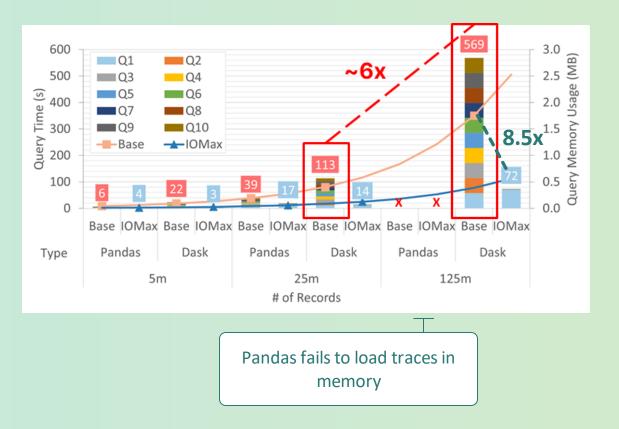
Workload	# of Records	# of Files	# of Processes	Size
1000 Genomes	715,248,240	21,268,291	2,712	440GB
Montage	12,346,353	19,680	11,488	30GB

- Hardware
 - Lassen supercomputer at LLNL
 - IBM Power9 CPU, 256GB RAM
 - IBM Spectrum Scale FS (GPFS)
- Software
 - Pandas (Data analysis library)
 - Utilized for in-memory analysis
 - Dask (Distributed computing library)
 - Utilized for out-of-core analysis
 - Pandas compatible APIs
 - Recorder (I/O tracing library)
 - Provides fine-grained I/O events

Evaluations \rightarrow Data Reduction



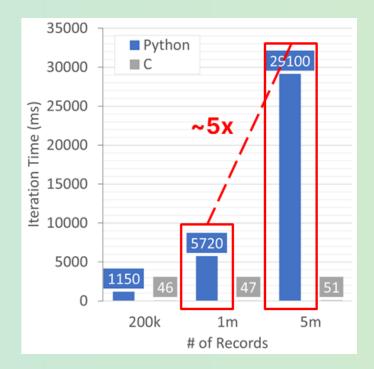
- 10 real-world I/O analysis queries are executed against datasets to showcase the effects of independent query treatment and the increasing memory footprint
- The execution time of the unoptimized queries increases linearly, ~6x per scale, while optimized version scales well
- The optimized version has a **memory footprint that is 8.5x smaller** than the unoptimized version



Evaluations \rightarrow Iterative Queries



- Performing I/O analysis and data drilling involves iterative operations. Typically, this is part of the analysis process. In our case, we offloaded these iterative operations to our preprocessor (DatasetTransformer), which is written in C
- The findings reveal a linear increase in Python's iteration time, ~5x per scale
- In contrast, C exhibits a constant iteration time, averaging around 50 ms

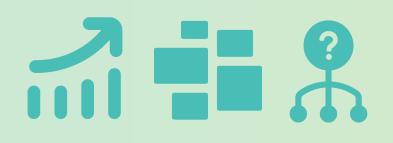


Evaluations → Datatype Performance

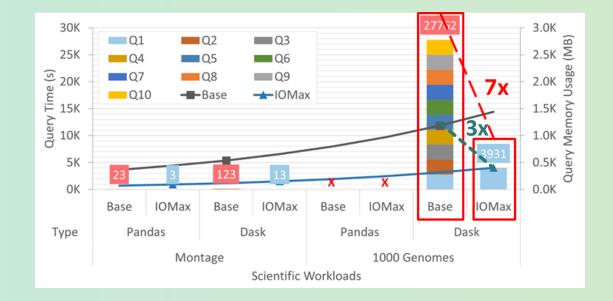
- Range queries were executed on both string and integer indices to identify the top ½, ¼, and ½ of the most accessed files
 - Our DataTransformer transforms string columns like filename into integer through hashing
- Accessing subsets with string indices exhibits a linear increase in access time relative to the number of files, resulting in an 11.4x slower performance
 - In contrast, integer indices shows a nearly constant access time, averaging around 15µs
- String indices incur a **memory footprint 13.1x larger** than that of integer indices



Evaluations → Scientific Workflows

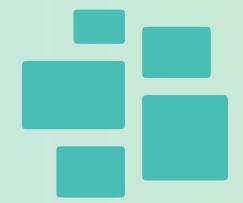


- 10 real-world I/O analysis queries were executed against the I/O traces of Montage and 1000 Genomes to illustrate the overall benefits derived from our methodology
- The finding show that IOMax **improves the analysis performance up to 7x** for large-scale I/O traces
- Additionally, it reduces the memory footprint of queries by 3x



Conclusion







Analyzing terabyte-scale I/O traces

Our methodology improves real-world large-scale I/O analysis performance by 7x Inefficient slicing and grouping

Our data transformer **improves** data-slicing performance by 11.4x Lack of global query optimization

Our query optimizations achieve up to 8.6x performance improvement and an 11x reduction in memory usage

Thank you!

Any questions?

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