

Using The Active Storage Fabrics Model To Address Petascale Storage Challenges

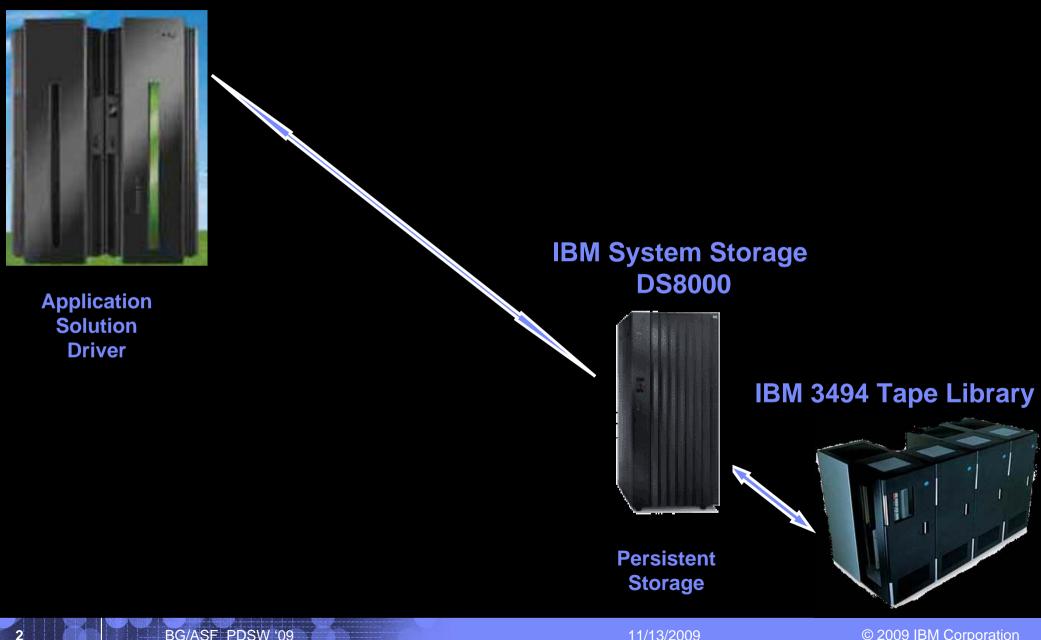
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4th Petascale Data Storage Workshop Sunday November 15, 2009 Portland Convention Center

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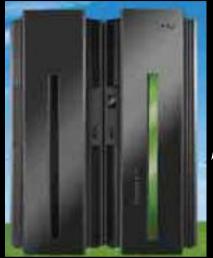


BG/ASF starts with a standard enterprise data center... **IBM Enterprise Server**





...and adds parallel processing in the storage hierarchy. IBM Enterprise Server



Application Solution Driver



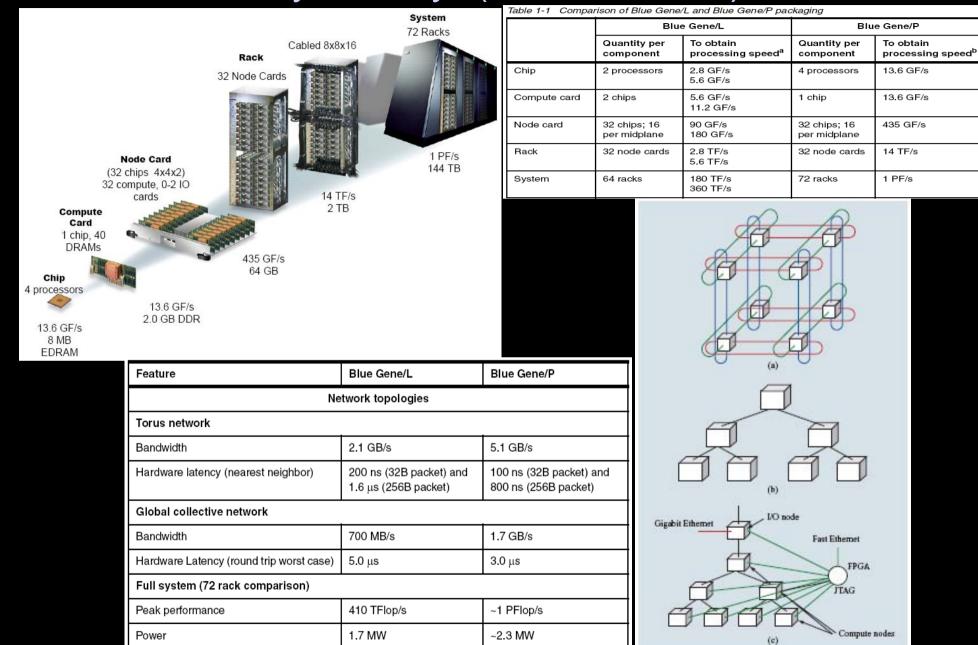
Managed as scalable Solid State Storage with Embedded Process Modules

Persistent Storage



11/13/2009

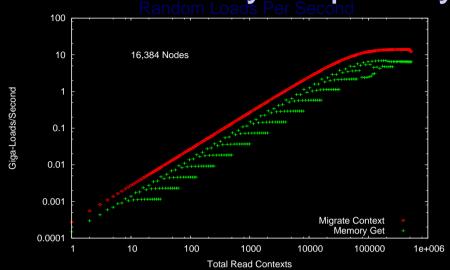
Blue Gene Family History (BG/L & BG/P)

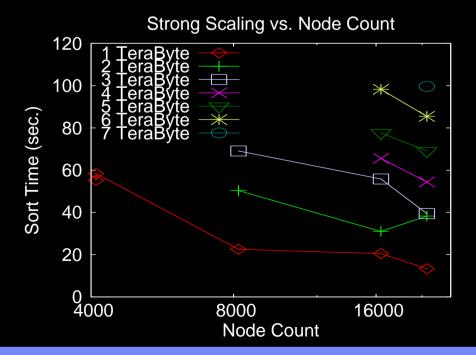


11/13/2009

Blue Gene Hardware Processor-In-Memory Capability

- Blue Gene/L micro benchmarks show order of magnitude improvement over largest SMP and clusters
- Scalable random pointer chasing through terabytes of memory
 - BG/L 1.0 GigaLoads / Rack
 - 14 GigaLoads/s for 16 racks
 - P690 0.3 Gigaloads
- Scalable "Indy" sort of terabytes in memory
 - BG/L 1.0 TB sort in 10 secs
 - BG/L 7.5 TB sort in 100 secs
 - Linux cluster, 80 nodes, 2530 disks 435 seconds (disk-to-disk)
- Streams benchmark
 - BG/L 2.5TB / Sec / Rack
 - P5 595 0.2TB / Sec / Rack





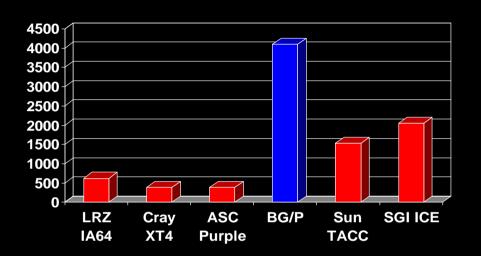
BG/ASF PDSW '09

11/13/2009

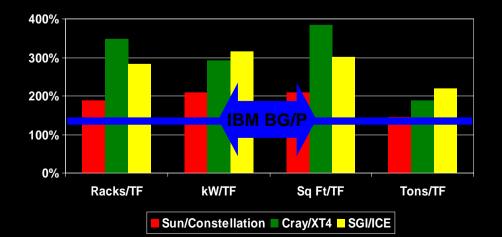
Main Memory Capacity per Rack



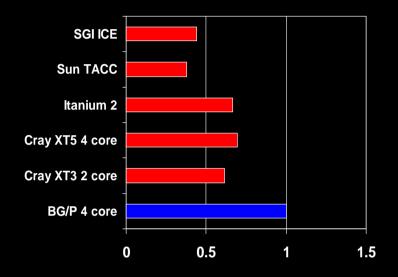
Blue Gene System-On-A-Chip Advantage



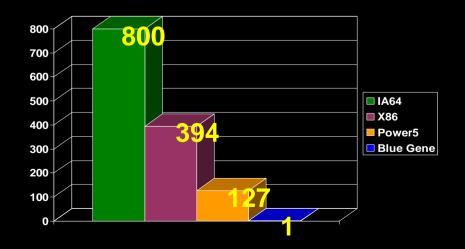
Relative power, space and cooling efficiencies



Peak Memory Bandwidth per node (byte/flop)

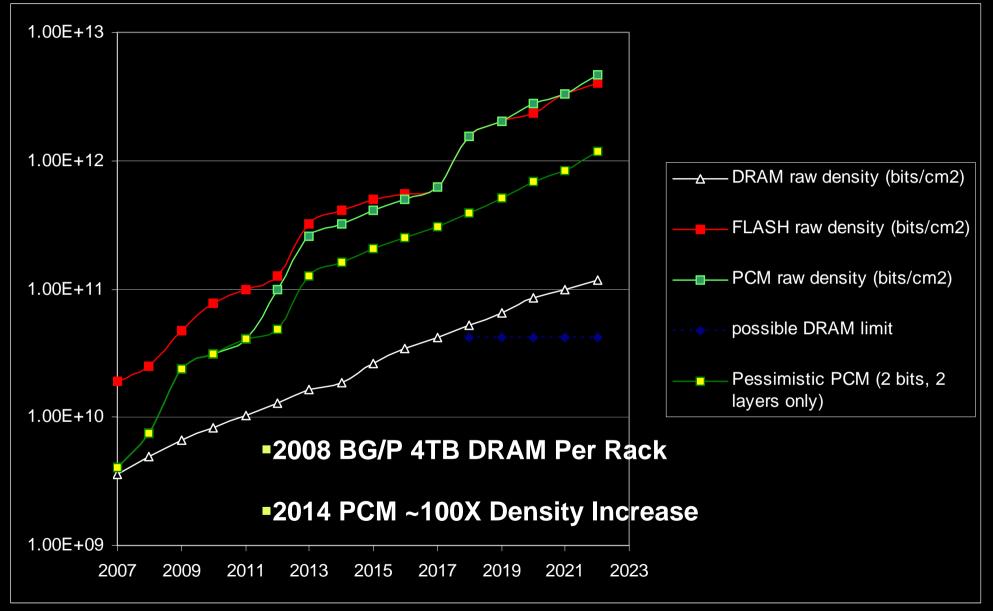


Failures per Month per @ 100 TF





Phase Change Memory (PCM) Density Timeline



ITRS 2008 data

Scalable, System-on-a-Chip, Storage Class Memory Platform

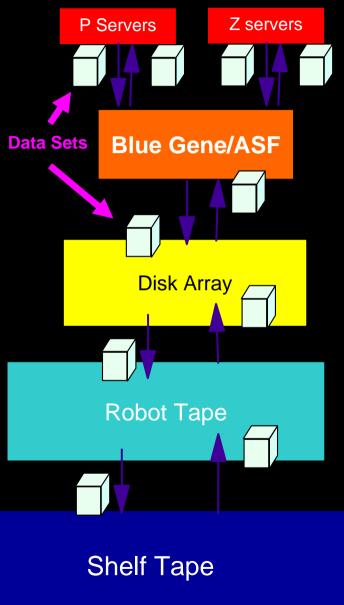
- Hypothetical SCM BG/ASF system:
 - − 1024 BG nodes, each with 1TB of SCM → 1 Petabyte Capacity
 - SCM bandwidth to be balanced with network all-to-all capability
 - bounded by bisectional bandwidth
 - Forms a persistent active storage fabric

BG/ASF allows hosts to transparently exploit large storage class memories

- Offload of bottleneck operations avoids pulling data through the main processor complex
- resiliency model reduces need to transfer data outside of fabric
- Current Blue Gene systems already make available sufficient memory for SCM systems research

Blue Gene / Active Storage Fabrics (BG/ASF) Concept

- Manage Blue Gene hardware as a scalable, solid state storage device with embedded processing capability
- Integrate this "Active Storage Fabric" with middleware such as DB2, MySql, GPFS, etc, at the data/storage interface using a parallel, key-value in-memory data store
- Transparently accelerate server Jobs with ASF:
 - Create libraries of reusable, low overhead
 Embedded Parallel Modules (EPMs) (scan, sort, join, sum, etc)
 - EPMs directly access middleware data (tables, files, etc) in key/value DB based on middleware data model
 - Middleware makes ASF datasets available via legacy interfaces allowing incremental parallelization and interoperability with legacy middleware function and host programs
- Integrate BG/ASF with information life-cycle and other standard IT data management products





BG/ASF Acceleration Opportunities

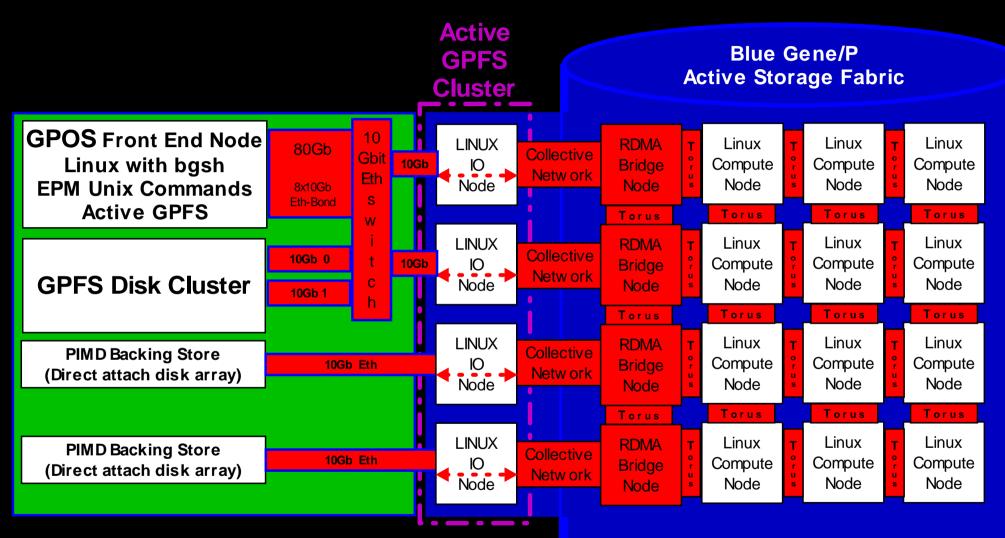
Embarrassingly Parallel

- Table scan
- Batch serial
 - Floating point intensive
 - Job farm/cloud
- Other applications without internal global data dependencies

Network Dependent

- Join
- Fast sort
 - "order by", "group by" queries
- Map-Reduce
- Aggregation operations
 - count(), sum(), min(), max(), avg(), ...
- Data analysis/OLAP
 - Aggregation with "group by"...
 - Real-time analytics
- HPC applications

BG/ASF : Accelerating Unix Utilities in Active GPFS



Blue Gene/L ASF: Early Performance Comparison with Host Commands

Benchmark process*:

- Generate a 22 GB synthetic "Indy Benchmark" data file
- grep for records with "AAAAAA" (outputs 1/3'd of the data which is used as input to sort)
- sort the output of grep

* NOTE: All ASF data taken on untuned system, non-optimized and with full debug printfs, all p55 data taken with unix time on command line)

Function	BG/ASF	pSeries p55
(Measured at Host shell and break out of EPM components)	512 nodes	(1 thread)
Total host command time: create Indy File:	(22GB) 22.26 s	686s
Total host command time: unix grep of Indy File:	25.1s	3180s
File Blocks to PIMD Records (Embedded grep) :	7.0 s	
Grep core function (Embedded grep) :	2.5s	
PIMD Records to File Blocks (Embedded grep) :	7.6s	
Total time on Blue Gene/ASF fabric (Embedded grep) :	18.8s	
Total host command time: sort output of grep:	60.41s	2220s
File Blocks to PIMD Records (Embedded sort) :	2.9s	
Sort core function (Embedded sort) :	12.2s	
PIMD Records to File Blocks (Embedded sort) :	15.8s	
Total time on Blue Gene/ASF fabric (Embedded sort) :	54.7s	



Blue Gene/L ASF: Early Scalability Results

Benchmark process*:

- Generate a 22 GB synthetic "Indy Benchmark" data file
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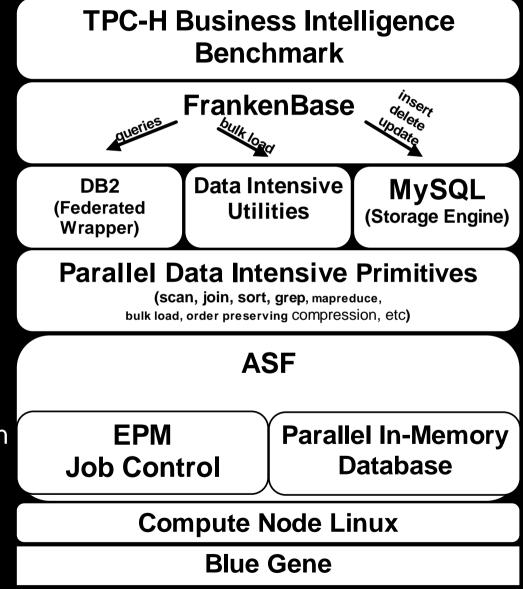
Function	BG/ASF	BG/ASF
	512 nodes	8K nodes
Total host command time: create Indy File:	(22GB) 22.26 s	(1TB) 197s
Total host command time: unix grep of Indy File:	25.1s	107s
File Blocks to PIMD Records (Embedded grep) :	7.0 s	18.1s
Grep core function (Embedded grep) :	2.5s	5.1s
PIMD Records to File Blocks (Embedded grep) :	7.6s	19.9s
Total time on Blue Gene/ASF fabric (Embedded grep) :	18.8s	50.2s
Total host command time: sort output of grep:	60.41s	120s
File Blocks to PIMD Records (Embedded sort) :	2.9s	6.8s
Sort core function (Embedded sort) :	12.2s	14.8s
PIMD Records to File Blocks (Embedded sort) :	15.8s	22.4s
Total time on Blue Gene/ASF fabric (Embedded sort) :	54.7s	66.55s

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BG/ASF Relational Database Acceleration

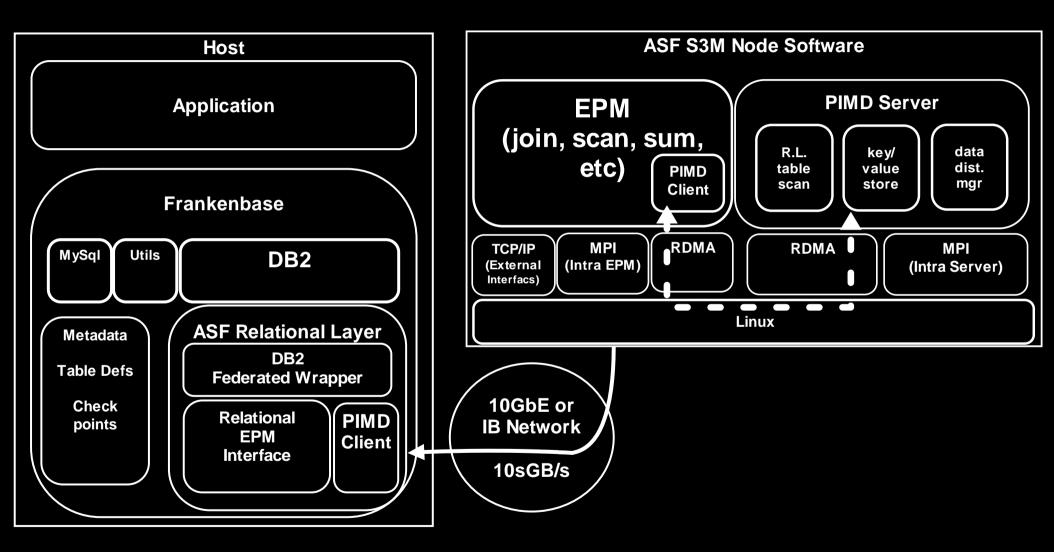
- Explore accelerating RDBMS systems by integrating at the data/storage interface
 - MySQL as an active storage engine
 - Defines tables, does inserts/updates
 - DB2 using Infosphere Federation
 - Offloads queries including scans, joins
 - Some exploitation of DB2 optimizer
- ASF Relational Layer adds
 - External table definition
 - Row scan capability to PIMD using query embedded field information – streaming, parallel, predicate evaluation
 - Embedded parallel join built using MPI and row scan
 - Potential for aggregation

BG/ASF PDSW '09





Frankenbase Component Diagram



Use case: Shared Molecular Dynamics Scientific Warehouse

- Very large data volumes which will increase with HPC system capability – expensive to move and store relative to the cost of production
- Datasets are often reused for many investigations beyond the initial inquiry – a mix of ad-hoc analytics are employed
- Currently, analytic capability and collaborative opportunities can be limited by the ability to access massive molecular dynamics datasets and the resources needed to perform ad-hoc queries
- Use of the ASF approach to implement a shared, active, data warehouse is expected to increase the scientific impact of a future super computer facilities

Research Community

Molecular Dynamics Analysis

Analytics

Active RDBMS Active File System

Active Storage Fabric

S3M



Summary

Current Status

- Demonstrated ASF accelerating unix utilities in Active GPFS
- Built BG/P ASF node environment (Linux, MPICH, RDMA, PIMD, etc)
- Completed initial function for ASF based DB2 acceleration
 - ASF Relational Layer
 - Tuple management system for storing table rows in key/value store
 - Query specific compilation for tuple predicates, projections and joins
 - PIMD server based row scan service supporting predicates and projections
 - Parallel join
 - DB2 Federated wrapper allows offload of DB2 selects (scans, joins)
 - MySQL to parse CREATE TABLE and produce tuple metadata
 - TPC-H table generation
 - PIMD Check point / reload infrastructure



Backup



Active Storage Fabrics Key Enablers

- Parallel In-Memory Database (PIMD)
 - The shared data storage model

- Embedded Parallel Modules (EPMs)
 - The storage embedded parallel programming model

Parallel In Memory Database (PIMD)

- Parallel In Memory Database Overview
 - Key/Value database
 - Client/Server architecture
 - Clients can be parallel programs running on the same nodes as servers
 - ... or on external nodes
 - Multiple datasets are cataloged by name (called "Partitioned Data Sets" (PDSs))
 - Multiuser support with UNIX-like PDS access management
 - Clients access records via a target server node determined by a hash function.
 - Server group retains control over actual data locality but optimizes for hash target
- Requirements/Challenges
 - 10k-100k nodes each with 1/P of the data (1GB...16GB per node)
 - Manage distribution including dynamic redistribution of data
 - Minimize overheads esp. IPC between Client<->Server and Internal (MPI)
- Future research
 - Fault tolerance for platforms beyond Blue Gene/P
 - Lock management and transactions
 - Persistence
 - Security



Embedded Parallel Module (EPMs)

- Parallel programs that use the Parallel In Memory Database (PIMD)
 - Currently fairly normal MPI programs
 - Use the PIMD Client interface to open PIMD Partitioned Data Sets (PDS).
 - Input and Output operands should be in PIMD PDSs
 - Use locality aware PIMD iterators to work on records that are on or near the node the EPM task is running on.
 - Worst case, an application driven data decomposition is an all-to-all away

Current Runtime Environment

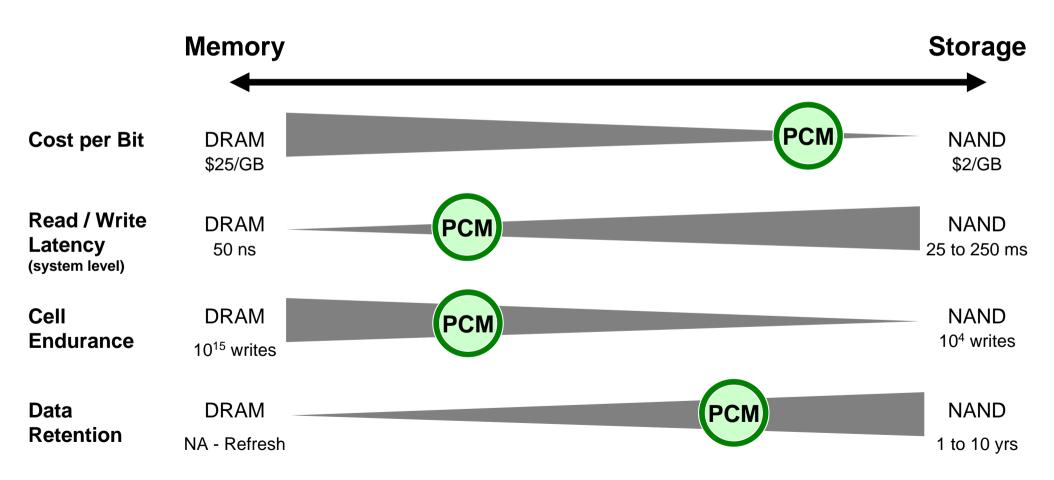
- Linux on Blue Gene/P compute nodes EPMs and the PIMD Server group share the same BG/P nodes.
- Client <-> Server communications via standard iWARP RDMA (soon, OFED verbs)
- Client internal communications via MPICH over TCP/IP on Blue Gene/P Networks
- EPM job Scheduling currently mpirun; expect to move to Load Leveler soon
- EPMs encourage parallel solution modularity
 - Exchange operands via PIMD an analog to UNIX file system and pipes
 - Initial EPMs will include a subset of UNIX text utilities
- Future research
 - Can be executed in pre-constructed MPI partitions already internally connected and to connected to PIMD Servers (These we call "Virtual Partitions")

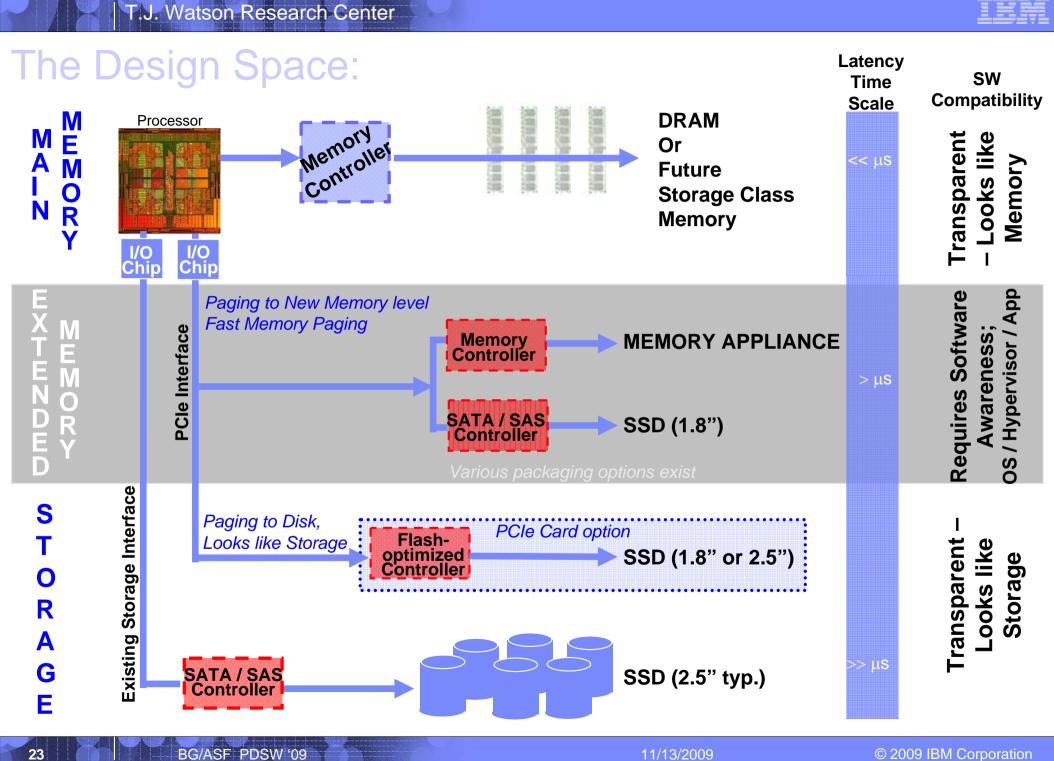


The Application Landscape

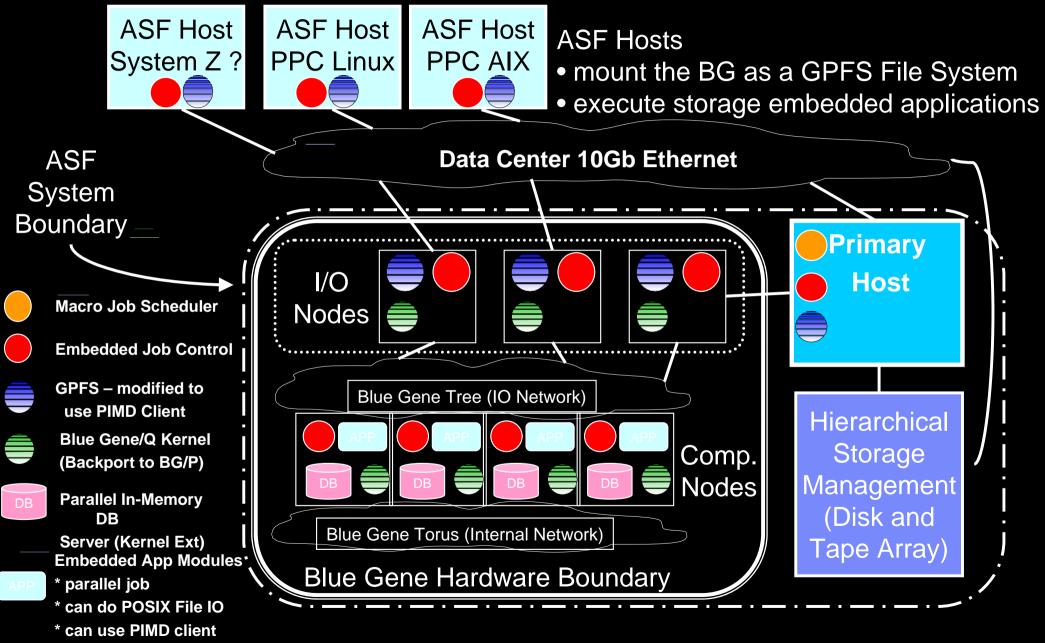
What can be done with a material with many of the advantages of DRAM but projected cost and retention closer to Flash?

Note: The numbers shown below are to act as reference points, not absolute targets/requirements.



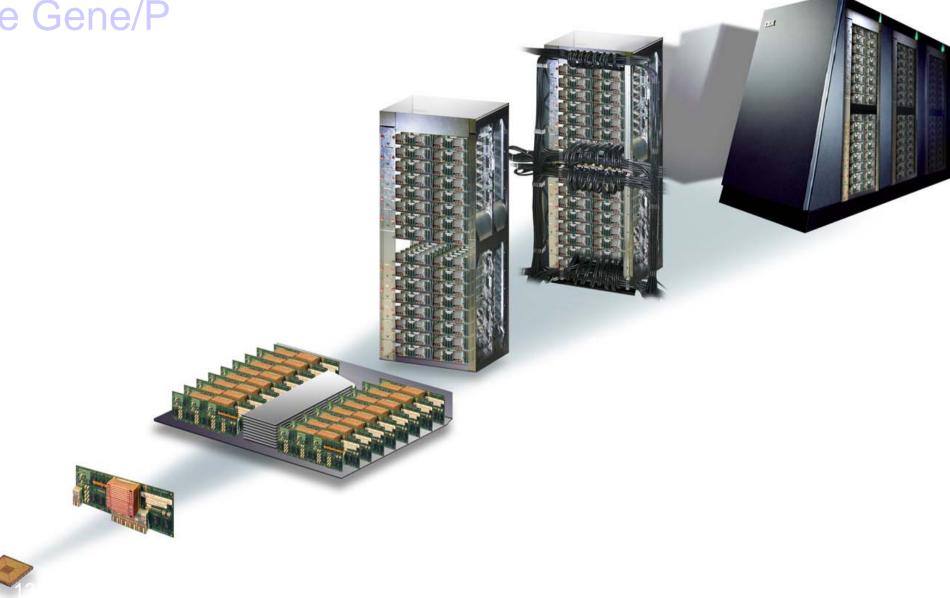


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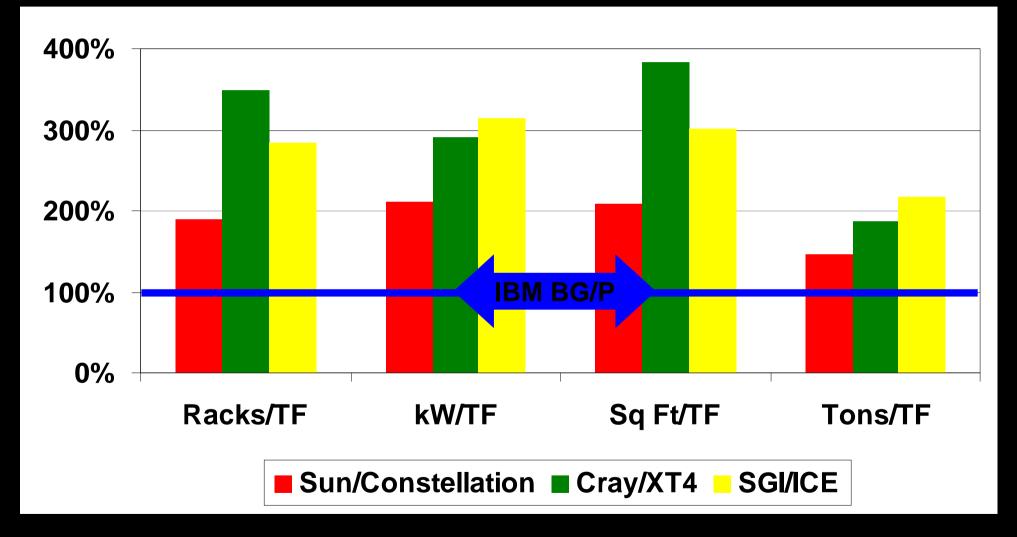
System

Blue Gene/P





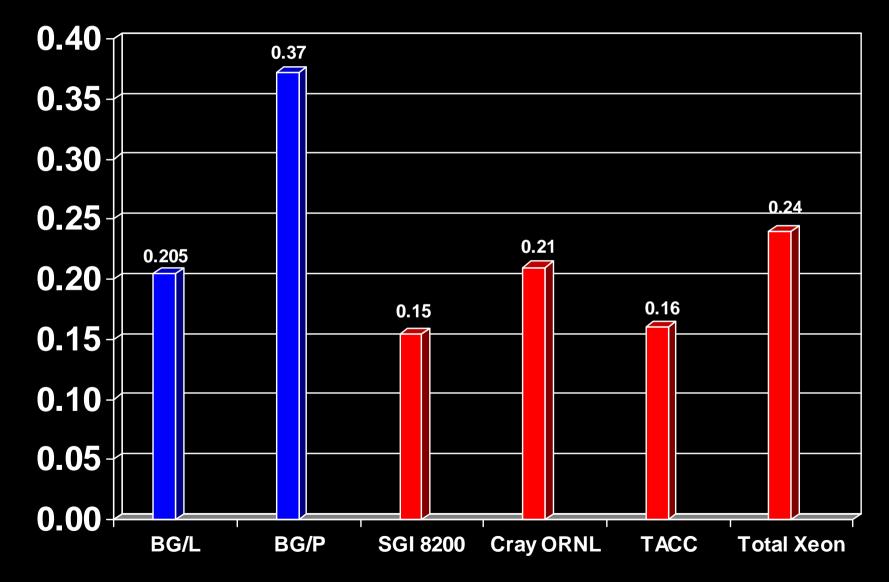
Relative power, space and cooling efficiencies (Published specs per peak performance)



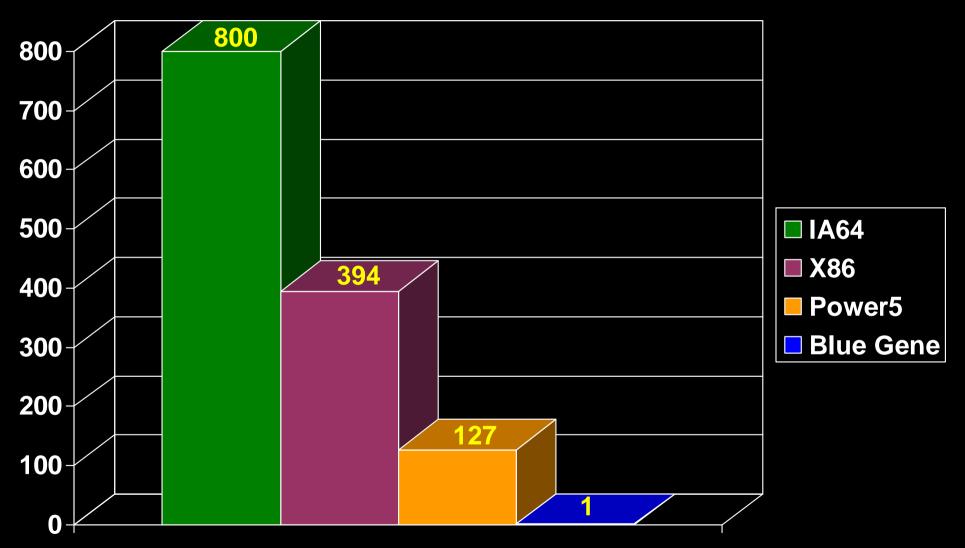


Green 500





Failures per Month per @ 100 TFlops (20 BG/L racks) unparalleled reliability



Results of survey conducted by Argonne National Lab on 10 clusters ranging from 1.2 to 365 TFlops (peak); excluding storage subsystem, management nodes, SAN network equipment, software outages



Blue Gene Technology Roadmap

