

Lawrence Livermore National Laboratory

Pianola: A script-based I/O benchmark



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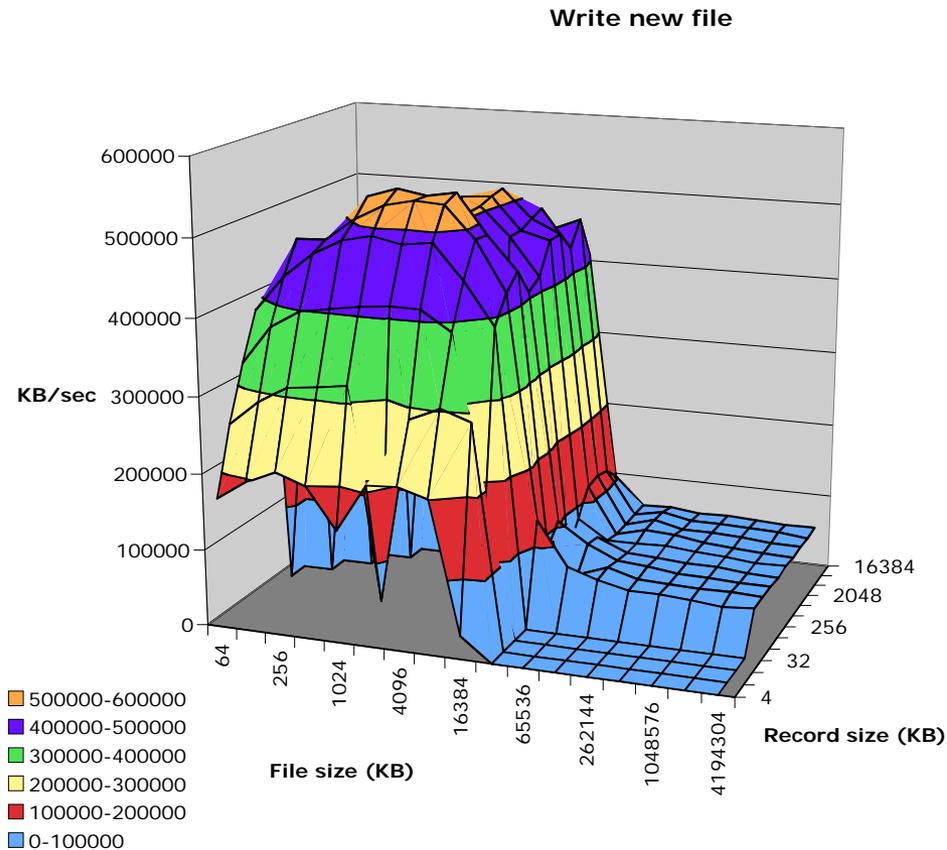
I/O benchmarking: What's going on here?

- Is my computer's I/O system "fast"?
- Is the I/O system keeping up with my application?
- Is the app using the I/O system effectively?
- What tools do I need to answer these questions?

- And what exactly do I mean by "I/O system" anyway?
 - For this talk, an I/O system is everything involved in storing data from the filesystem to the storage hardware



Existing tools can measure general or application-specific performance



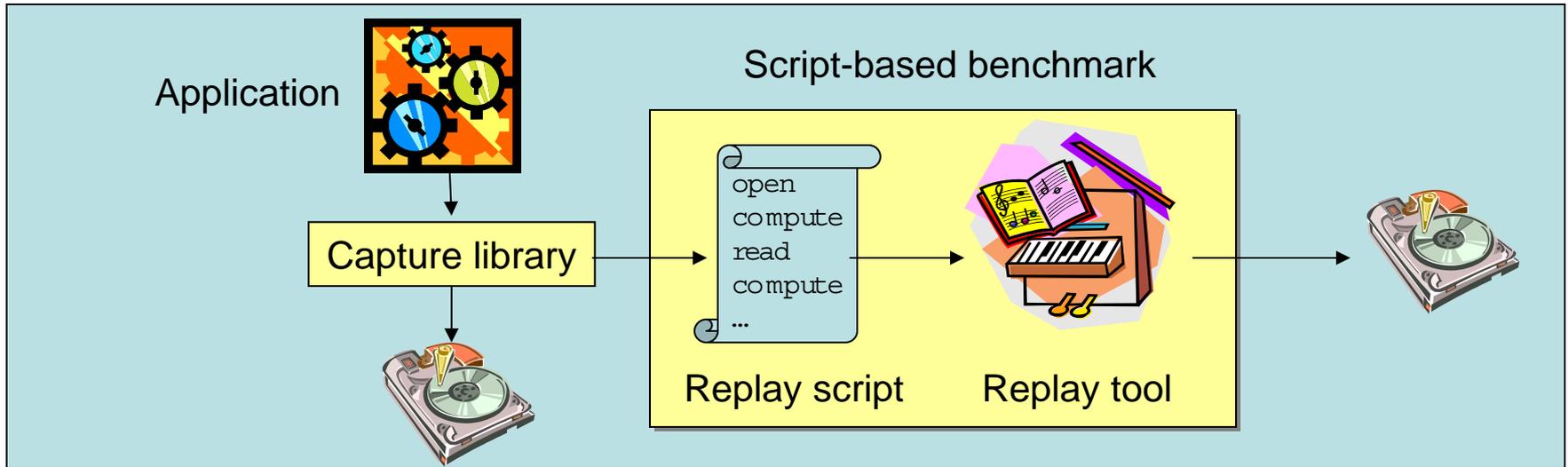
- IOzone automatically measures I/O system performance for different operations and parameters
- Relatively little ability to customize I/O requests
- Many application-oriented benchmarks
 - SWarp, MADbench2...
- Interface-specific benchmarks
 - IOR, //TRACE

Measuring the performance that matters

- System benchmarks only measure general response, not application-specific response
- Third-party application-based benchmarks may not generate the stimulus you care about
- In-house applications may not be practical benchmarks
 - Difficult for nonexperts to build and run
 - Nonpublic source cannot be distributed to vendors and collaborators
- Need benchmarks that...
 - Can be generated and used easily
 - Model application-specific characteristics

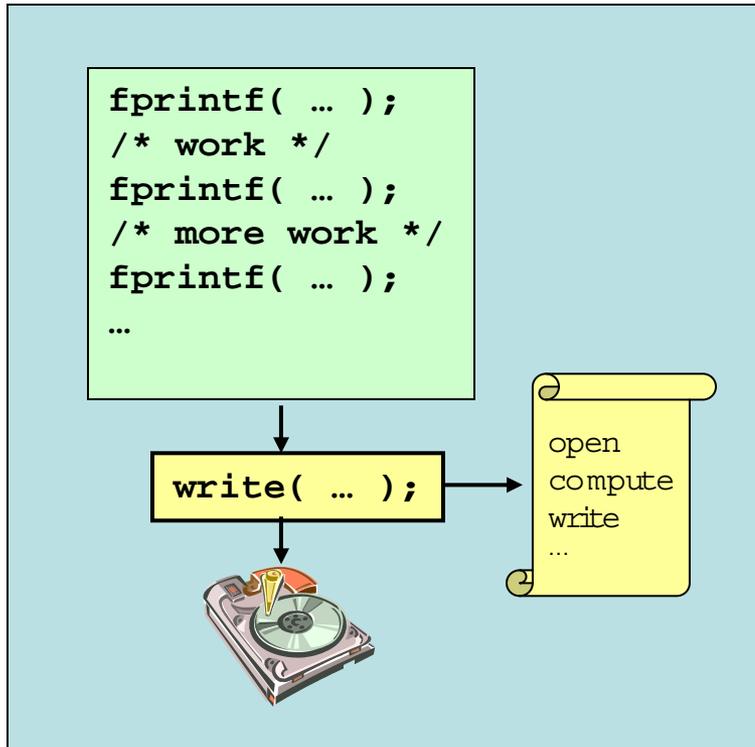


Script-based benchmarks emulate real apps



- Capture trace data from application and generate the same sequence of operations in a replay-benchmark
- We began with //TRACE from CMU (Ganger's group)
 - Records I/O events and intervening "compute" times
 - Focused on parallel I/O, but much of the infrastructure is useful for our *sequential* I/O work

Challenges for script-based benchmarking: Recording I/O calls at the the right level



- Instrumenting at high level
 - + Easy with LD_PRELOAD
 - Typically generates more events, so logs are bigger
 - Need to replicate formatting
 - Timing includes computation
- Instrumenting at low level
 - + Fewer types of calls to capture
 - + Instrumentation is at I/O system interface
 - Cannot use LD_PRELOAD to intercept all calls

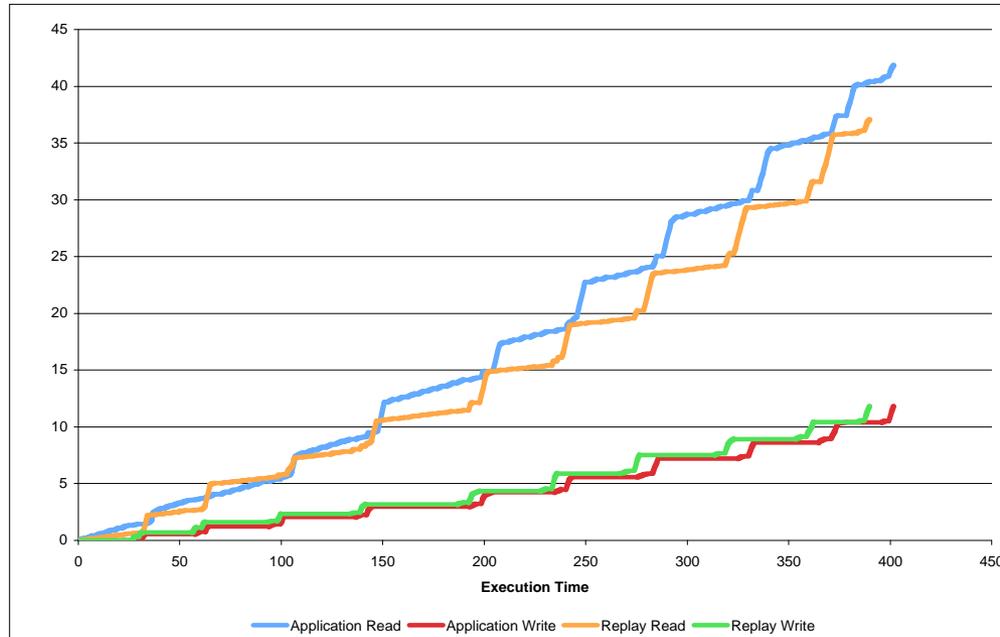
First attempt at capturing system calls: Linux strace utility

```
$ strace -r -T -s 0 -e trace=file,desc ls
0.000000 execve("/bin/ls", [ ..], [/* 29 vars */]) = 0 <0.000237>
0.000297 open("/etc/ld.so.cache", O_RDONLY) = 3 <0.000047>
0.000257 fstat64(3, {st_mode=S_IFREG|0644, st_size=64677, ...}) = 0 <0.000033>
0.000394 close(3) = 0 <0.000015>
0.000230 open("/lib/libc.so.1", O_RDONLY) = 3 <0.000046>
0.000289 read(3, ""..., 512) = 512 <0.000028>
...
```

- Records any selected set of system calls
- Easy to use: just add to command line
- Produces easily parsed output



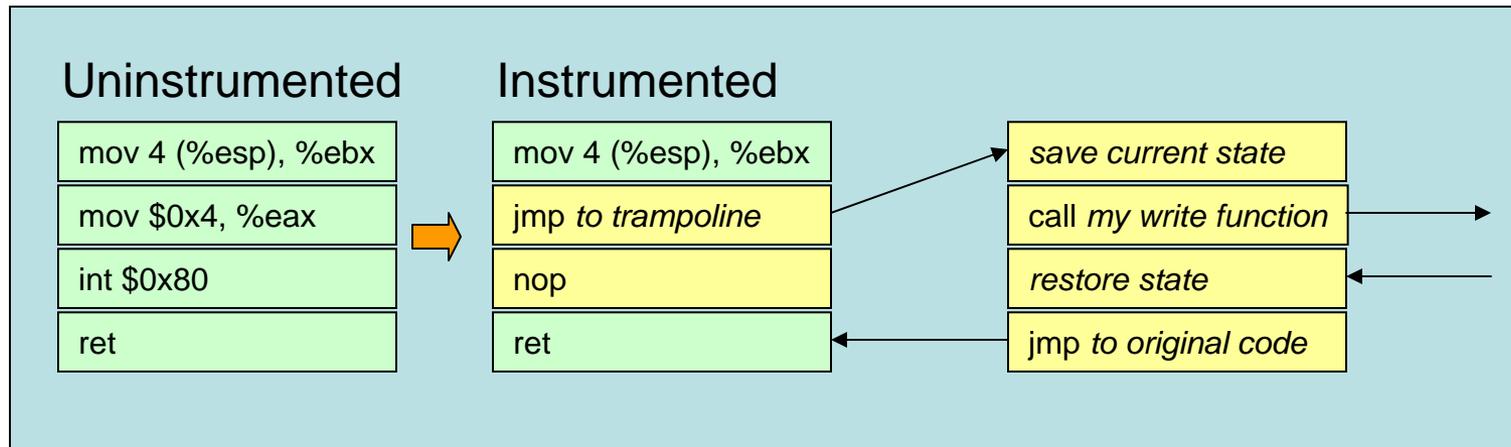
Strace results look reasonably accurate, but overall runtime is exaggerated



	Read (sec.)	Write (sec.)	Elapsed (sec.)
Uninstrumented	--	--	324
Instrumented Application	41.8	11.8	402
Replay	37.0	11.8	390



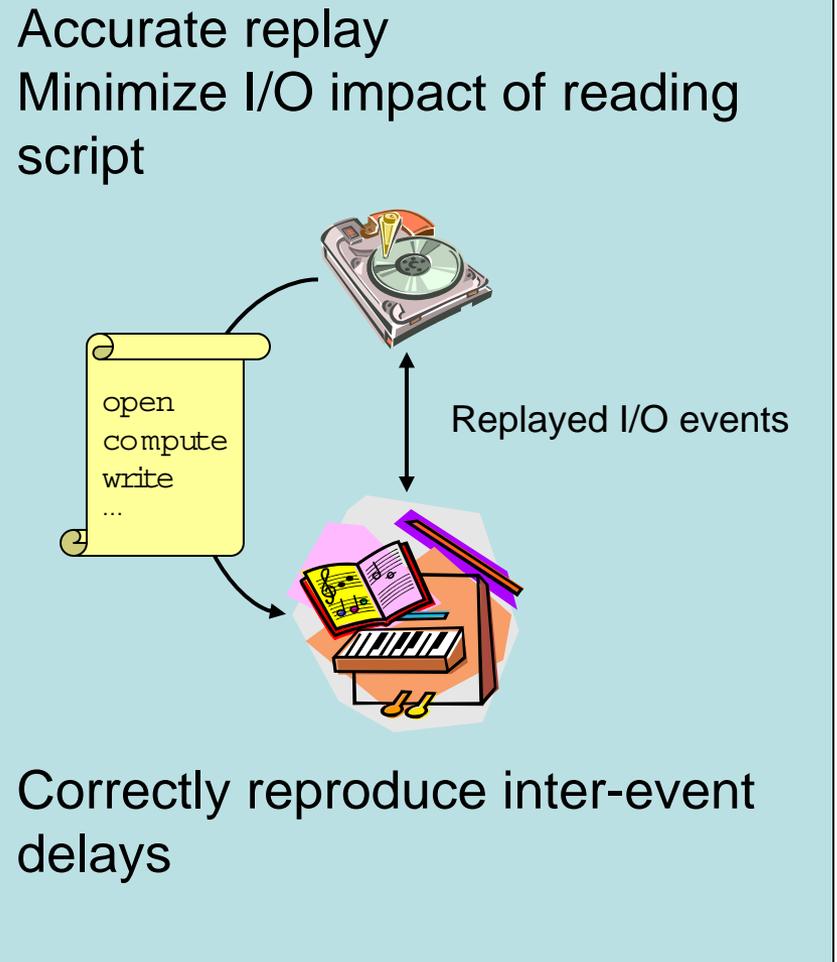
For accurate recording, gather I/O calls using binary instrumentation



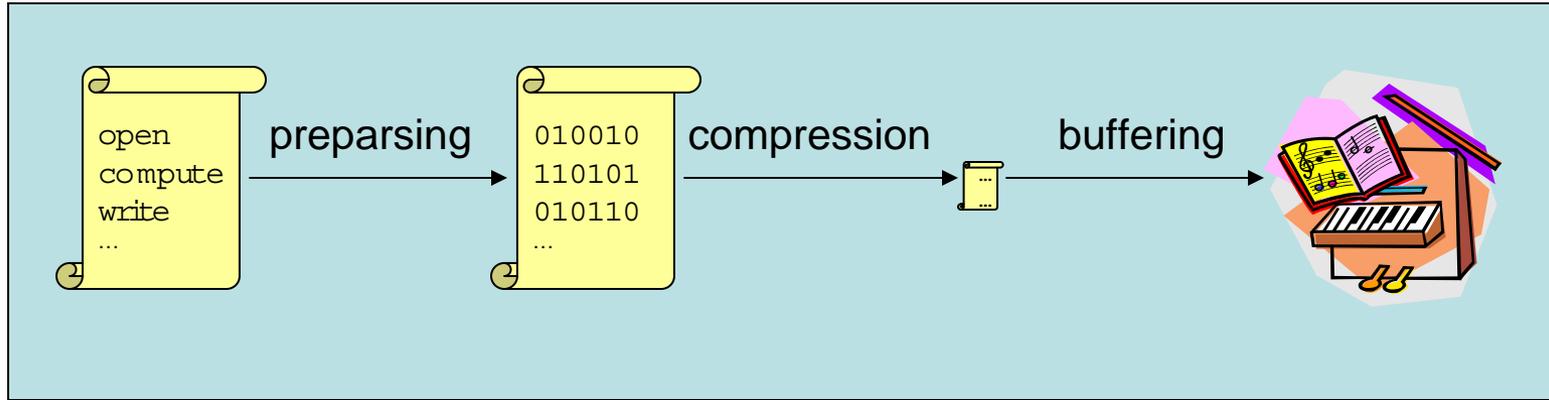
- Can intercept and instrument specific system-level calls
- Overhead of instrumentation is paid at program startup
- Existing Jockey library works well for x86_32, but not ported to other platforms
- Replay can be portable, though

Issues for accurate replay

- Replay engine must be able to read and parse events quickly
- Reading script must not interfere significantly with I/O activities being replicated
- Script must be portable across platforms

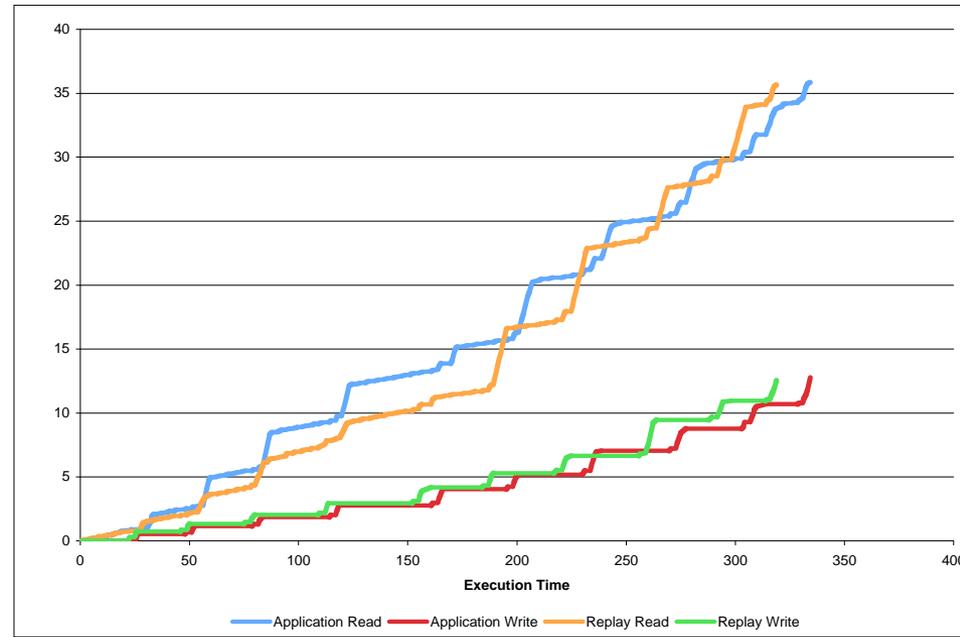


Accurate replay: Preparsing, compression, and buffering



- Text-formatted output script is portable across platforms
- Instrumentation output is parsed into binary format and compressed (~30:1)
 - Conversion done on target platform
- Replay engine reads and buffers script data during “compute” phases between I/O events

Replay timing and profile match original application well



	Read (sec.)	Write (sec.)	Elapsed(sec.)
Uninstrumented	--	--	314
Instrumented Application	35.8	12.8	334
Replay	35.7	12.5	319



Things that didn't help

- Compressing text script as it's generated
 - Only 2:1 compression
 - Time of I/O events themselves are not what's very important during instrumentation phase
- Replicating the memory footprint
 - Memory used by application is taken from same pool as I/O buffer cache
 - Smaller application (like the replay engine) should go faster because more buffer space available
 - Replicated memory footprint by tracking `brk()` and `mmap()` calls, but it made no difference!



Conclusions on script-based I/O benchmarking

- Gathering accurate I/O traces is harder than it seems
 - Currently, no solution is both portable and efficient
- Replay is easier, but efficiency still matters
- Many possibilities for future work—which matter most?
 - File name transformation
 - Parallel trace and replay
 - More portable instrumentation
 - How to monitor mmap'd I/O?

