

Implementation, evaluation and analysis of Block index for ADIOS

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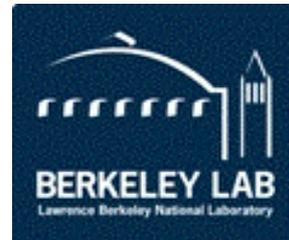


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Introduction

- Scientific datasets are commonly stored and managed by parallel file systems and I/O libraries
 - E.g. Lustre, HDF5, NetCDF, ADIOS
 - optimized for reading/writing **large chunks of data**
 - Data layout and file organization impact query performance
- The characteristics and behaviors of I/O systems should be considered into the design of indexing methods

The idea of “Block index”

- Indexing blocks (consecutive data records) instead of individual data records
 - Reduce index size
 - Reduce number of I/O requests
 - Reading an individual record has similar I/O latency as reading a data block

Implement block index into ADIOS

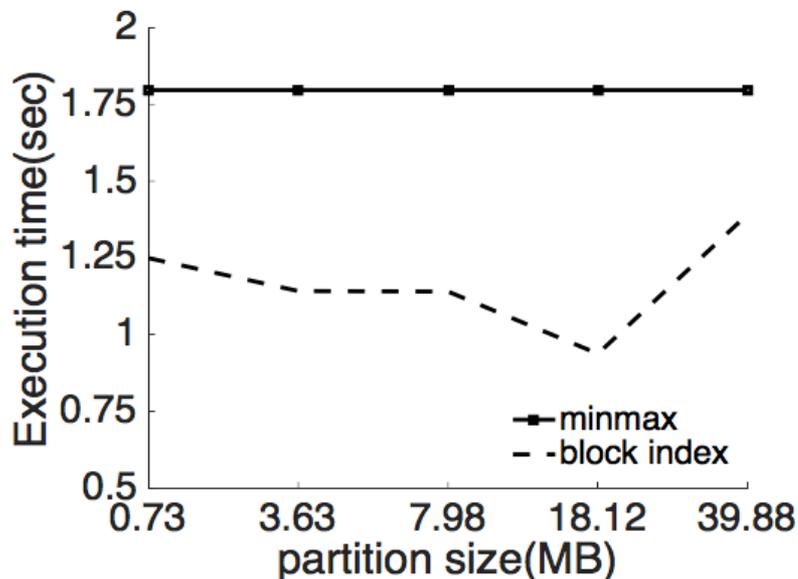
- Minmax method in ADIOS
 - Records the min, max value from each writeblock
 - The size of writeblock => the size of data of each process (can be extremely big)
- Block index method in ADIOS
 - **Logically** divides a writeblock into smaller partitions
 - Records the min, max values of each partition
 - Using logical partition can maintain the same number of writeblock
 - The I/O requests on the same writeblock can be merged by ADIOS to minimize I/O contention

Experiment Setup

- Edison Cray XC30 at NERSC
 - 5576 compute nodes, with 12-core Intel Ivy Bridge 2.4GHz CPU and 64GB memory per node
 - Lustre parallel file system with 72GB peak performance
- S3D dataset
 - Each variable contains $1100 \times 1080 \times 1408$ double precision records
 - Each variable is written to file using 64 writeblocks of size $275 \times 270 \times 352$ (~200MB)

Performance evaluation

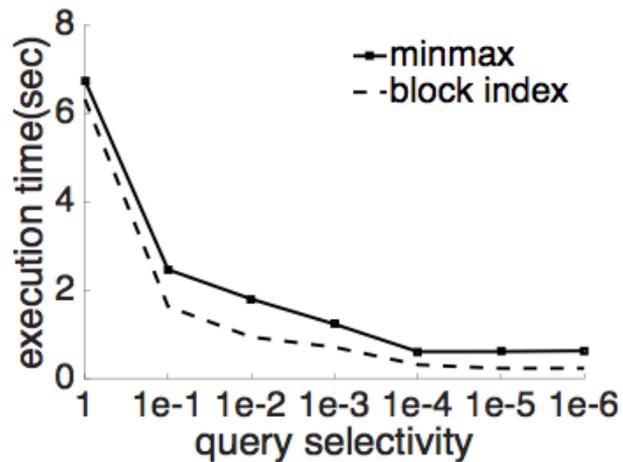
- Varied partition size
 - The performance is a tradeoff between read size and I/O throughput
 - Minmax's read bytes is more than twice the block index



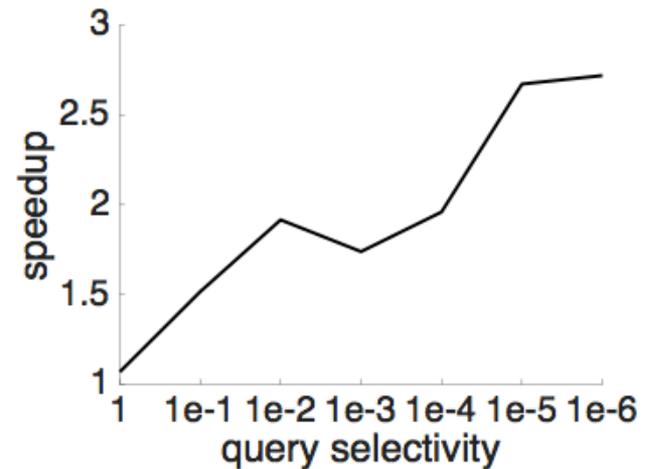
| Partition size | read requests | bytes read | I/O throughput |
|----------------|---------------|------------|----------------|
| 0.73MB | 1298 | 941.17MB | 753.76MB/s |
| 3.63MB | 266 | 964.38MB | 852.84MB/s |
| 7.98MB | 124 | 989.03MB | 867.29MB/s |
| 18.12MB | 59 | 1069.52MB | 1141.22MB/s |
| 39.88MB | 30 | 1196.41MB | 864.65MB/s |
| minmax | 11 | 2193.42MB | 1222.17MB/s |

Performance evaluation

- Varied query selectivity
 - Block index reads less data when query selectivity is smaller => speedup is higher
 - Similar performance under 100% query selectivity



(a) Execution time.



(b) speedup of block index.

Conclusion

- Query performance of minmax is limited by the size of writeblock
- Query performance of Block index that logically partitions a writeblock improves due to less data reading, and more flexible read size
- Future work
 - Performance analysis and modeling of I/O systems
 - Design the algorithm to select the proper block size and request merging condition



THANK YOU